

## What is a Graph?

- A graph consists of two sets  $V$  and  $E$  where

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$$V = \{v_1, \dots, v_n\}$$

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$$E = \{e_1, \dots, e_m\}$$

- where each edge  $e_i$  consists of an **unordered** vertex pair  $e_i = \{v_j, v_k\}$ .
- $j \neq k$  because we avoid self loops.
- For a *directed* graph we consider edges as ordered pairs.

## Example

- Let  $V = \{1, 2, 3, 4, 5\}$  and
- $E = \{(1, 2), (1, 3), (2, 3), (2, 4)\}$ .
- For an unordered graph, the order of the vertices in each edge does not matter.
- $G = \{V, E\}$  is a graph.

## Drawing a graph

Let us draw

## Complete Graphs

- A complete graph means that every vertex is connected to every other vertex.
- In other words for all  $i, j$  we have  $e_k = \{v_i, v_j\} \in E$ .
- We avoid self loops so  $i \neq j$ .
- We denote them by  $K_n$ .

## Graph Vertex Coloring

- For a graph  $G = \{V, E\}$  a *vertex coloring* is a choice of colors
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$$C = \{c_1, \dots, c_p\}$$

- and an assignment of colors to vertices  $f : V \rightarrow C$
- such that for any edge  $e = \{v_i, v_j\}$ ,
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$$f(v_i) \neq f(v_j).$$

- In other words, for each edge we assign its constituent vertices with different colors.

## Example

### Question: Chromatic Number

- Given a graph  $G = \{V, E\}$ ,
- what is the least number of colors to use to vertex color the graph?
- We denote the number by  $\chi(G)$ .
- It is called the *chromatic number* for the graph.

### Bounds on Chromatic Number

- Claim:  $\chi(G) \leq n$  where  $n$  is number of vertices of  $G$ .
- Proof:
- For complete graphs equality holds,  $\chi(G) = n$ .
- Can we find a better bound?

### contd..

- Let us look at the degrees at each vertex.
- Here *degree* of a vertex is the number of edges emanating from the vertex.
- Example:

### contd..

- Suppose  $\Delta(G)$  is the largest degree of any vertex.
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$$\chi(G) \leq \Delta(G) + 1.$$

- Can we show this?
- Answer: Use a greedy algorithm.

## Application: Class Scheduling

- If there are  $n$  classes.
- Students sign up for them without knowing the time schedules.
- How many distinct time slots to pick?
- Answer:

## Answer:

- Pick vertices as each class.
- Join two vertices by an edge if there is a student taking both classes.
- Chromatic number gives the minimum number of distinct time slots.

## LP Formulation

- Can you formulate an LP for the above problem?
- Say we have a graph  $G = (V, E)$ .
- To start off we can pick  $n$  colors  $\{c_1, \dots, c_n\}$ .
- Variables:
- $y[i] = 1$  if we use color  $i$  and 0 otherwise.
- $x[i, j] = 1$  if vertex  $i$  is assigned color  $j$ , 0 otherwise

## Chromatic LP

- Minimize  $\sum_i y[i]$ .
- subject to
- $\sum_k x[ik] = 1$  for every  $i$ .
- $x[ik] \leq y[k]$
- $x[ik] + x[jk] \leq 1$  for  $(v_i, v_k) \in E$  and each  $k$ .
- $y[k] \leq y[k - 1]$ .
- $x[ik], y[k] \in \{0, 1\}$ .

## Edge Coloring

- Question: How many colors do we need to color the edges of a graph such that no two edges of the same color connect at a vertex?
- Example:

## Bounds on Edge Coloring

- We need at least  $\Delta(G)$  colors.
- Infact this is a very good estimate.

- Vizing in 1964 proved that the number of colors needed is either
- $\Delta(G)$  or
- $\Delta(G) + 1$ .

## Edge Coloring Example

- Schedule matches in a tournament.
- Suppose we have 9 teams and we have a graph saying which teams must play each other.
- How many rounds of play needed to achieve this?

## Example:

## LP Formulation

In [ ]: